## Features

- Utilizes the $A V R^{\circledR}$ Enhanced RISC Architecture
- 120 Powerful Instructions - Most Single Clock Cycle Execution
- 4K bytes of In-System Reprogrammable Downloadable Flash
- SPI Serial Interface for Program Downloading
- Endurance: 1,000 Write/Erase Cycles
- 256 bytes EEPROM
- Endurance: 100,000 Write/Erase Cycles
- 256 bytes Internal SRAM
- $32 \times 8$ General Purpose Working Registers
- 32 Programmable I/O Lines
- Programmable Serial UART
- SPI Serial Interface
- $\mathrm{V}_{\mathrm{cc}}$ : 2.7-6.0V
- Fully Static Operation, 0-20 MHz
- Instruction Cycle Time: 50 ns @ 20 MHz
- One 8-Bit Timer/Counter with Separate Prescaler
- One 16-Bit Timer/Counter with Separate Prescaler and Compare and Capture Modes
- Dual PWM
- External and Internal Interrupt Sources
- Programmable Watchdog Timer with On-Chip Oscillator
- On-Chip Analog Comparator
- Low Power Idle and Power Down Modes
- Programming Lock for Software Security


## Description

The AT90S4414 is a low-power CMOS 8-bit microcontroller based on the AVR ${ }^{\circledR}$ enhanced RISC architecture. By executing powerful instructions in a single clock cycle, the AT90S4414 achieves throughputs approaching 1 MIPS per MHz allowing the system designer to optimize power consumption versus processing speed.
The AVR core combines a rich instruction set with 32 general purpose working registers. All the 32 registers are directly connected to the Arithmetic Logic Unit (ALU), allowing two independent registers to be accessed in one single instruction executed in one clock cycle. The resulting architecture is more code efficient while achieving throughputs up to ten times faster than conventional CISC microcontrollers.

## Pin Configurations




8-Bit $A V \boldsymbol{R}^{\text {® }}$ Microcontroller with 4K bytes Downloadable Flash

AT90S4414 Preliminary

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## Block Diagram



Figure 1. The AT90S4414 Block Diagram

## Description (Continued)

The AT90S4414 provides the following features: 4 K bytes of downloadable Flash, 256 bytes EEPROM, 256 bytes SRAM, 32 general purpose I/O lines, 32 general purpose working registers, flexible timer/counters with compare modes, internal and external interrupts, a programmable serial UART, programmable watchdog timer with internal oscillator, an SPI serial port and two software selectable power saving modes. The idle mode stops the CPU while allowing the SRAM, timer/ counters, SPI port and interrupt system to continue functioning. The power down mode saves the register contents but freezes the oscillator, disabling all other chip functions until the next interrupt or hardware reset.
The device is manufactured using Atmel's high density non-volatile memory technology. The on-chip downloadable Flash allows the program memory to be reprogrammed in-system through an SPI serial interface or by a conventional nonvolatile memory programmer. By combining an enhanced RISC 8-bit CPU with downloadable Flash on a monolithic chip, the Atmel AT90S4414 is a powerful microcontroller that provides a highly flexible and cost effective solution to many embedded control applications.
The AT90S4414 AVR is supported with a full suite of program and system development tools including: C compilers, macro assemblers, program debugger/simulators, in-circuit emulators, and evaluation kits.

## Architectural Overview

The fast-access register file contains $32 \times 8$-bit general purpose working registers with a single clock cycle access time. This means that during one single clock cycle, one ALU (Arithmetic Logic Unit) operation is executed. Two operands are output from the register file, the operation is executed, and the result is stored back in the register file - in one clock cycle.
Six of the 32 registers can be used as three 16-bit indirect address register pointers for data space addressing - enabling efficient address calculations. One of the three address pointers is also used as the address pointer for the constant table look up function. These added function registers are the 16-bit X-register, Y-register and Z-register.
The ALU supports arithmetic and logic functions between registers or between a constant and a register. Single register operations are also executed in the ALU.
In addition to the register operation, conventional memory addressing modes can be used on the register file as well. This is enabled by the fact that the register file is assigned the 32 lowermost data space addresses ( $\$ 00-\$ 1 \mathrm{~F}$ ), allowing them to be accessed as though they were ordinary memory locations.
The I/O memory space contains 64 addresses for CPU peripheral functions as control registers, timer/counters, A/D-converters, and other I/O functions. The I/O Memory can be accessed directly, or as the Data Space locations following those of the register file, \$20-\$5F.
The AVR is a Harvard architecture - with separate memories and buses for program and data. The program memory is executed with single level pipelining. While one instruction is being executed, the next instruction is pre-fetched from the program memory. This concept enables instructions to be executed in every clock cycle. The program memory is in-system downloadable Flash memory.
With the relative jump and call instructions, the whole 4 K byte address space is directly accessed. Most AVR instructions have a single 16 -bit word format. Every program memory address contains a 16- or 32-bit instruction.
During interrupts and subroutine calls, the return address program counter (PC) is stored on the stack. The stack is effectively allocated in the general data SRAM, and consequently the stack size is only limited by the total SRAM size and the usage of the SRAM. All user programs must initialize the SP in the reset routine (before subroutines or interrupts are executed). The 16 -bit stack pointer SP is read/write accessible in the I/O space.
The 256 bytes data SRAM can be easily accessed through the five different addressing modes supported in the AVR architecture.

The memory spaces in the AVR architecture are all linear and regular memory maps.
A flexible interrupt module has its control registers in the I/O space with an additional global interrupt enable bit in the status register. Each interrupt has a separate interrupt vector in the interrupt vector table at the beginning of the program memory. Interrupts have priority in accordance with their interrupt vector position. The lower the interrupt address vector the higher priority.


Figure 2. Memory Maps

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## AT90S4414 Register Summary



AT90S4414 Instruction Set Summary

| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| ARITHMETIC AND LOGIC INSTRUCTIONS |  |  |  |  |  |
| ADD | Rd, Rr | Add two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}$ | Z,C,N,V,H | 1 |
| ADC | Rd, Rr | Add with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}+\mathrm{C}$ | Z,C,N,V,H | 1 |
| ADIW | Rdl, K | Add Immediate to Word | Rdh:Rdl $\leftarrow$ Rdh:Rdl + K | Z,C,N,V,S | 2 |
| SUB | Rd, Rr | Subtract two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}$ | Z,C,N,V,H | 1 |
| SUBI | Rd, K | Subtract Constant from Register | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}$ | Z,C,N,V,H | 1 |
| SBC | Rd, Rr | Subtract with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z,C,N,V,H | 1 |
| SBCI | Rd, K | Subtract with Carry Constant from Reg. | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}-\mathrm{C}$ | Z,C,N,V,H | 1 |
| SBIW | Rdl, K | Subtract Immediate from Word | Rdh:Rdl $\leftarrow$ Rdh:Rdl - K | Z,C,N,V,S | 2 |
| AND | Rd, Rr | Logical AND Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rr}$ | Z,N,V | 1 |
| ANDI | Rd, K | Logical AND Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{K}$ | Z,N,V | 1 |
| OR | Rd, Rr | Logical OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \vee \mathrm{Rr}$ | Z,N,V | 1 |
| ORI | Rd, K | Logical OR Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rd} v \mathrm{~K}$ | Z,N,V | 1 |
| EOR | Rd, Rr | Exclusive OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rr}$ | Z,N,V | 1 |
| COM | Rd | One's Complement | $\mathrm{Rd} \leftarrow$ \$ FF - Rd | Z,C,N,V | 1 |
| NEG | Rd | Two's Complement | $\mathrm{Rd} \leftarrow$ \$00-Rd | Z,C,N,V,H | 1 |
| SBR | Rd,K | Set Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} v \mathrm{~K}$ | Z,N,V | 1 |
| CBR | Rd,K | Clear Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet(\$ \mathrm{FF}-\mathrm{K})$ | Z,N,V | 1 |
| INC | Rd | Increment | $\mathrm{Rd} \leftarrow \mathrm{Rd}+1$ | Z,N,V | 1 |
| DEC | Rd | Decrement | $\mathrm{Rd} \leftarrow \mathrm{Rd}-1$ | Z,N,V | 1 |
| TST | Rd | Test for Zero or Minus | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rd}$ | Z,N,V | 1 |
| CLR | Rd | Clear Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rd}$ | Z,N,V | 1 |
| SER | Rd | Set Register | $\mathrm{Rd} \leftarrow$ \$ FF | None | 1 |
| BRANCH INSTRUCTIONS |  |  |  |  |  |
| RJMP | k | Relative Jump | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 2 |
| IJMP |  | Indirect Jump to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 2 |
| RCALL | k | Relative Subroutine Call | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 3 |
| ICALL |  | Indirect Call to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 3 |
| RET |  | Subroutine Return | $\mathrm{PC} \leftarrow$ STACK | None | 4 |
| RETI |  | Interrupt Return | $\mathrm{PC} \leftarrow$ STACK | I | 4 |
| CPSE | Rd, Rr | Compare, Skip if Equal | if (Rd = Rr) $\mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | $1 / 2$ |
| CP | Rd, Rr | Compare | Rd - Rr | Z, N,V,C,H | 1 |
| CPC | Rd, Rr | Compare with Carry | $\mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z, N,V,C,H | 1 |
| CPI | Rd,K | Compare Register with Immediate | Rd-K | Z, N,V,C,H | 1 |
| SBRC | $\mathrm{Rr}, \mathrm{b}$ | Skip if Bit in Register Cleared | if $(\operatorname{Rr}(\mathrm{b})=0) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | $1 / 2$ |
| SBRS | $\mathrm{Rr}, \mathrm{b}$ | Skip if Bit in Register is Set | if $(\operatorname{Rr}(\mathrm{b})=1) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | $1 / 2$ |
| SBIC | P, b | Skip if Bit in I/O Register Cleared | if $(\mathrm{P}(\mathrm{b})=0) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | $1 / 2$ |
| SBIS | P, b | Skip if Bit in I/O Register is Set | if $(\mathrm{P}(\mathrm{b})=1) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| BRBS | s, k | Branch if Status Flag Set | if (SREG(s) $=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRBC | s, k | Branch if Status Flag Cleared | if (SREG(s) $=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BREQ | k | Branch if Equal | if $(Z=1)$ then $P C \leftarrow P C+k+1$ | None | 1/2 |
| BRNE | k | Branch if Not Equal | if ( $Z=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRCS | k | Branch if Carry Set | if ( $\mathrm{C}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRCC | k | Branch if Carry Cleared | if ( $\mathrm{C}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRSH | k | Branch if Same or Higher | if ( $\mathrm{C}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRLO | k | Branch if Lower | if ( $\mathrm{C}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRMI | k | Branch if Minus | if $(\mathrm{N}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRPL | k | Branch if Plus | if ( $\mathrm{N}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRGE | k | Branch if Greater or Equal, Signed | if ( $\mathrm{N} \oplus \mathrm{V}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRLT | k | Branch if Less Than Zero, Signed | if ( $\mathrm{N} \oplus \mathrm{V}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRHS | k | Branch if Half Carry Flag Set | if ( $\mathrm{H}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRHC | k | Branch if Half Carry Flag Cleared | if ( $\mathrm{H}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRTS | k | Branch if T Flag Set | if ( $\mathrm{T}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRTC | k | Branch if T Flag Cleared | if $(T=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRVS | k | Branch if Overflow Flag is Set | if $(\mathrm{V}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRVC | k | Branch if Overflow Flag is Cleared | if $(\mathrm{V}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRIE | k | Branch if Interrupt Enabled | if $(\mathrm{I}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRID | k | Branch if Interrupt Disabled | if $(1=0)$ then $P C \leftarrow P C+k+1$ | None | $1 / 2$ |

DATA TRANSFER INSTRUCTIONS

| MOV | Rd, Rr | Move Between Registers | $\mathrm{Rd} \leftarrow \mathrm{Rr}$ | None | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| LDI | Rd, K | Load Immediate | $\mathrm{Rd} \leftarrow \mathrm{K}$ | None | 1 |
| LD | Rd, X | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{X})$ | None | 2 |
| LD | Rd, $\mathrm{X}_{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{X}), \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| LD | Rd, - X | Load Indirect and Pre-Dec. | $\mathrm{X} \leftarrow \mathrm{X}-1, \mathrm{Rd} \leftarrow(\mathrm{X})$ | None | 2 |
| LD | Rd, Y | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LD | Rd, Y+ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Y}), \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| LD | Rd, - Y | Load Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1, \mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LDD | Rd, $\mathrm{Y}+\mathrm{q}$ | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Y}+\mathrm{q})$ | None | 2 |
| LD | Rd, Z | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LD | Rd, $\mathrm{Z}_{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Z}), \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| LD | Rd, -Z | Load Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1, \mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LDD | Rd, $\mathrm{Z}+\mathrm{q}$ | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Z}+\mathrm{q})$ | None | 2 |
| LDS | Rd, k | Load Direct from SRAM | $\mathrm{Rd} \leftarrow(\mathrm{k})$ | None | 3 |
| ST | X, Rr | Store Indirect | $(\mathrm{X}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{X}+$, Rr | Store Indirect and Post-Inc. | $(\mathrm{X}) \leftarrow \mathrm{Rr}, \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| ST | - X, Rr | Store Indirect and Pre-Dec. | $X \leftarrow X-1,(X) \leftarrow \operatorname{Rr}$ | None | 2 |
| ST | Y, Rr | Store Indirect | $(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{Y}+$, Rr | Store Indirect and Post-Inc. | $(\mathrm{Y}) \leftarrow \mathrm{Rr}, \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| ST | - $\mathrm{Y}, \mathrm{Rr}$ | Store Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1,(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | $\mathrm{Y}+\mathrm{q}, \mathrm{Rr}$ | Store Indirect with Displacement | $(\mathrm{Y}+\mathrm{q}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | Z, Rr | Store Indirect | $(\mathrm{Z}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | Z+, Rr | Store Indirect and Post-Inc. | $(\mathrm{Z}) \leftarrow \mathrm{Rr}, \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| ST | -Z, Rr | Store Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1,(\mathrm{Z}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | $\mathrm{Z}+\mathrm{q}, \mathrm{Rr}$ | Store Indirect with Displacement | $(\mathrm{Z}+\mathrm{q}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STS | k, Rr | Store Direct to SRAM | $(\mathrm{k}) \leftarrow \mathrm{Rr}$ | None | 3 |
| LPM |  | Load Program Memory | $\mathrm{R} 0 \leftarrow(\mathrm{Z})$ | None | 3 |
| IN | Rd, P | In Port | $\mathrm{Rd} \leftarrow \mathrm{P}$ | None | 1 |
| OUT | $\mathrm{P}, \mathrm{Rr}$ | Out Port | $\mathrm{P} \leftarrow \mathrm{Rr}$ | None | 1 |
| PUSH | Rr | Push Register on Stack | STACK $\leftarrow \mathrm{Rr}$ | None | 2 |
| POP | Rd | Pop Register from Stack | $\mathrm{Rd} \leftarrow$ STACK | None | 2 |
| BIT AND BIT-TEST INSTRUCTIONS |  |  |  |  |  |
| SBI | P, b | Set Bit in I/O Register | $\mathrm{l} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 1$ | None | 2 |
| CBI | P, b | Clear Bit in I/O Register | $\mathrm{l} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 0$ | None | 2 |
| LSL | Rd | Logical Shift Left | $\mathrm{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \operatorname{Rd}(0) \leftarrow 0$ | Z,C,N,V | 1 |
| LSR | Rd | Logical Shift Right | $\operatorname{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \operatorname{Rd}(7) \leftarrow 0$ | Z,C,N,V | 1 |
| ROL | Rd | Rotate Left Through Carry | $\mathrm{Rd}(0) \leftarrow \mathrm{C}, \mathrm{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \mathrm{C} \leftarrow \operatorname{Rd}(7)$ | Z,C,N,V | 1 |
| ROR | Rd | Rotate Right Through Carry | $\mathrm{Rd}(7) \leftarrow \mathrm{C}, \mathrm{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \mathrm{C} \leftarrow \operatorname{Rd}(0)$ | Z,C,N,V | 1 |
| ASR | Rd | Arithmetic Shift Right | $\mathrm{Rd}(\mathrm{n}) \leftarrow \mathrm{Rd}(\mathrm{n}+1), \mathrm{n}=0 . .6$ | Z,C,N,V | 1 |
| SWAP | Rd | Swap Nibbles | $\operatorname{Rd}(3 . .0) \leftarrow \operatorname{Rd}(7 . .4), \operatorname{Rd}(7 . .4) \leftarrow \operatorname{Rd}(3 . .0)$ | None | 1 |
| BSET | s | Flag Set | SREG(s) $\leftarrow 1$ | SREG(s) | 1 |
| BCLR | S | Flag Clear | SREG(s) $\leftarrow 0$ | SREG(s) | 1 |
| BST | Rr, b | Bit Store from Register to T | $\mathrm{T} \leftarrow \operatorname{Rr}(\mathrm{b})$ | T | 1 |
| BLD | Rd, b | Bit load from T to Register | $\operatorname{Rd}(\mathrm{b}) \leftarrow \mathrm{T}$ | None | 1 |
| SEC |  | Set Carry | $\mathrm{C} \leftarrow 1$ | C | 1 |
| CLC |  | Clear Carry | $\mathrm{C} \leftarrow 0$ | C | 1 |
| SEN |  | Set Negative Flag | $N \leftarrow 1$ | N | 1 |
| CLN |  | Clear Negative Flag | $\mathrm{N} \leftarrow 0$ | N | 1 |
| SEZ |  | Set Zero Flag | $\mathrm{Z} \leftarrow 1$ | Z | 1 |
| CLZ |  | Clear Zero Flag | $\mathrm{Z} \leftarrow 0$ | Z | 1 |
| SEI |  | Global Interrupt Enable | $1 \leftarrow 1$ | 1 | 1 |
| CLI |  | Global Interrupt Disable | $1 \leftarrow 0$ | 1 | 1 |
| SES |  | Set Signed Test Flag | $\mathrm{S} \leftarrow 1$ | S | 1 |
| CLS |  | Clear Signed Test Flag | $\mathrm{S} \leftarrow 0$ | S | 1 |
| SEV |  | Set Twos Complement Overflow. | $\mathrm{V} \leftarrow 1$ | V | 1 |
| CLV |  | Clear Twos Complement Overflow | $\mathrm{V} \leftarrow 0$ | V | 1 |
| SET |  | Set T in SREG | $\mathrm{T} \leftarrow 1$ | T | 1 |
| CLT |  | Clear T in SREG | $\mathrm{T} \leftarrow 0$ | T | 1 |
| SEH |  | Set Half Carry Flag in SREG | $\mathrm{H} \leftarrow 1$ | H | 1 |
| CLH |  | Clear Half Carry Flag in SREG | $\mathrm{H} \leftarrow 0$ | H | 1 |
| NOP |  | No Operation |  | None | 1 |
| SLEEP |  | Sleep | (see specific descr. for Sleep function) | None | 3 |
| WDR |  | Watchdog Reset | (see specific descr. for WDR/timer) | None | 1 |

